

# No-Free-Lunch Theorem for ML in Isabelle/HOL

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# Today's Talk

## Topic

A formalization of the no-free-lunch theorem for ML in Isabelle/HOL.

- Outline of the theorem
- Formalization in Isabelle/HOL

ML ... Machine Learning

## Reference

[1] *Understanding Machine Learning: From Theory to Algorithms*, Shai Shalev-Shwartz and Shai Ben-David, Cambridge University Press, 2014.

## Source Code

AFP entry: [\*No-free-lunch theorem for machine learning\*](#)

The views expressed are my own and do not necessarily reflect the views of my employer.

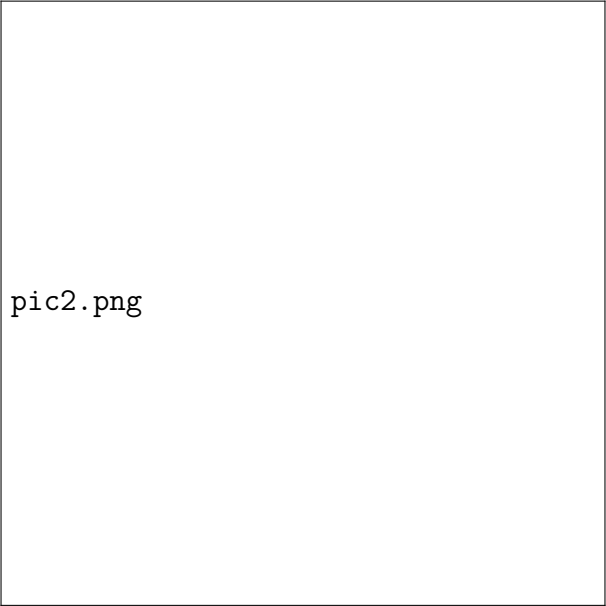
# No Free Lunch

TANSTAAFL is a popular adage.

*There ain't no such thing as a free lunch*

Wikipedia: [No such thing as a free lunch](#)

pic1\_1.png



pic2.png

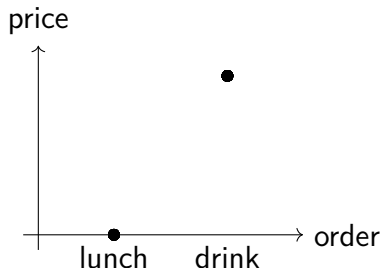
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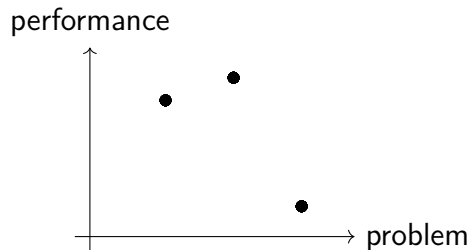
# No-Free-Lunch Theorem

It is impossible to get something for nothing.

**Free lunch restaurants**



**ML algorithms**



## No-Free-Lunch theorem for ML (the version in [1])

No universal learner exists.

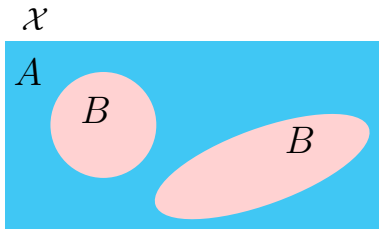
More precisely, for binary classification prediction tasks, for every learner, there exists a distribution on which it fails.



# Binary Classification

Binary classification ...

the task of putting each  $x \in \mathcal{X}$  into one of two categories.



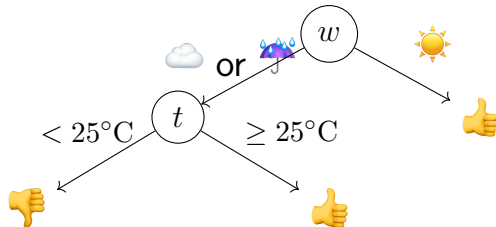
## Example

Predict whether ice cream sells well or not from weather conditions.

Domain set:  $\mathcal{X} = \text{Weather} \times \text{Temp}$

e.g.  $(\text{☀️}, 20^\circ\text{C}) \in \mathcal{X}$

Categories: group 👍 and group 👎



# Setting for Binary Classification Prediction Tasks

## Setting

Domain set:  $\mathcal{X}$

Label set:  $\mathbb{B} = \{0, 1\}$

Training data:  $S = ((x_1, b_1), \dots, (x_n, b_n))$ , where  $x_i \in \mathcal{X}$  and  $b_i \in \mathbb{B}$

Learning algorithm:  $A$ , where  $A(S) : \mathcal{X} \rightarrow \mathbb{B}$  is a predictor.

## Example

Predict whether ice cream sells well or not from weather conditions.

Domain set: Weather  $\times$  Temp

Training data:  $S = (((\text{☀️}, 20^\circ\text{C}), \text{👍}), ((\text{☔️}, 22^\circ\text{C}), \text{👎}), \dots)$

$A(S)$  might return the decision tree in the previous page. (e.g.  $A(S)(\text{☔️}, 27) = \text{👍}$ )

# The Method to Analyze Learning Algorithms

Assume that each training data is generated from some probability distribution.

i.e.,

$$(x, b) \sim \mathcal{D},$$

$S \sim \mathcal{D}^m$ , where  $\mathcal{D}$  is a probability distribution on  $\mathcal{X} \times \mathbb{B}$ .

## Error of predictor

$$\mathcal{L}_{\mathcal{D}}(h) = \Pr_{(x,b) \sim \mathcal{D}}[h(x) \neq b] \text{ for } h : \mathcal{X} \rightarrow \mathbb{B}.$$

## Bias of training data

is taken into account by assuming  $S \sim \mathcal{D}^m$ .

A *universal learning algorithm*  $A$  may has the following property. (cf. PAC learnability)

For all  $\varepsilon > 0$ ,  $\delta > 0$ , there exists  $M$  s.t. for all  $m \geq M$ , for all  $\mathcal{D}$ ,

$$\Pr_{S \sim \mathcal{D}^m} [\mathcal{L}_{\mathcal{D}}(A(S)) \leq \varepsilon] \geq 1 - \delta$$

# No-Free-Lunch Theorem

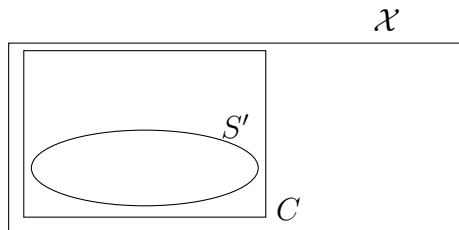
## No-Free-Lunch Theorem for ML

Let  $A$  be a learning algorithm over a domain  $\mathcal{X}$  and  $m < |\mathcal{X}|/2$ .  
Then, there exists a distribution  $\mathcal{D}$  on  $\mathcal{X} \times \mathbb{B}$  s.t.

- there exists  $f : \mathcal{X} \rightarrow \mathbb{B}$  s.t.  $\mathcal{L}_{\mathcal{D}}(f) = 0$ , and
- $\Pr_{S \sim \mathcal{D}^m} [\mathcal{L}_{\mathcal{D}}(A(S)) > 1/8] \geq 1/7$

Intuition:

Let  $C \subseteq \mathcal{X}$  s.t.  $|C| = 2m$ .  
Then, any learning algorithm  
has information only half of the  
instances in  $C$ .

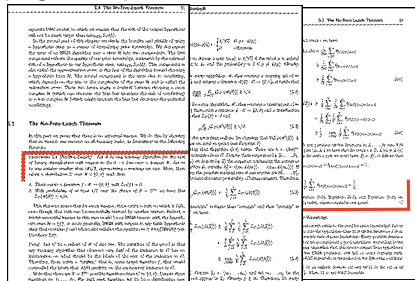


# Proof Overview

1. Define  $2^{|C|}$  uniform discrete distributions on  $C \times \mathbb{B}$ .
2. Show  $\max_D \mathbb{E}_{S \sim D^m} [\mathcal{L}_D(A(S))] \geq 1/4$ .
3. Apply the Markov inequality.

In the book  
2 pages of proof

In Isabelle/HOL  
435 lines of proof



```
theorem no_free_lunch_ML:
  fixes X :: "'a measure" and m :: nat
  and A :: "(nat  $\Rightarrow$  'a  $\times$  bool)  $\Rightarrow$  'a  $\Rightarrow$  bool"
  assumes X1:"finite (space X)  $\Rightarrow$  2 * m < card (space X)"
  and X2[measurable]:" $\lambda$ x. x  $\in$  space X  $\Rightarrow$  {x}  $\in$  sets X"
  and m[arith]:"0 < m"
  and A[measurable]: "( $\lambda$ (s,x). A s x)  $\in$  (PiM {.. $\leq$ m} ( $\lambda$ i. X
     $\rightarrow$  M count_space
  shows " $\exists D ::$  ('a  $\times$  bool) measure. sets D = sets (X  $\otimes$  M c
    prob_space D  $\wedge$ 
    ( $\exists$ f. f  $\in$  X  $\rightarrow$  M count_space (UNIV :: bool set))  $\wedge$ 
    P(s in Pi_M {.. $\leq$ m} ( $\lambda$ i. D). P((x, y) in D. A s
  proof -
    let D = "count_space (UNIV :: bool set)"
```

# Formalization in Isabelle/HOL

The type of learning algorithms in Isabelle/HOL.

$$A :: \underbrace{\text{nat}}_{\substack{\star \text{ the size} \\ \text{of data}}} \Rightarrow \left( \frac{\text{nat} \Rightarrow 'a \times \text{bool}}{S \in (\mathcal{X} \times \mathbb{B})^n} \right) \Rightarrow \frac{'a \Rightarrow \text{bool}}{\text{predictor}}$$

$\star$  is typically used to determine the number of loop executions.

We omit  $\star$  because we do not need the concrete definitions of algorithms.

**theorem** *no-free-lunch-ML*:

**fixes**  $X :: 'a \text{ measure}$  **and**  $m :: \text{nat}$

**and**  $A :: (\text{nat} \Rightarrow 'a \times \text{bool}) \Rightarrow 'a \Rightarrow \text{bool}$

**assumes**  $0 < m$  **and** *finite* (*space*  $X$ )  $\implies 2 * m < \text{card}(\text{space } X)$

**and**  $\bigwedge x. x \in \text{space } X \implies \{x\} \in \text{sets } X$  ← used to construct discrete distributions

**and**  $(\lambda(s,x). A \ s \ x) \in (\Pi_M i \in \{..<m\}. X \otimes_M \mathbb{B}) \otimes_M X \rightarrow_M \mathbb{B}$  ← measurability

**shows**  $\exists \mathcal{D} :: ('a \times \text{bool}) \text{ measure}. \text{sets } \mathcal{D} = \text{sets } (X \otimes_M \mathbb{B}) \wedge \text{prob-space } \mathcal{D} \wedge$

$(\exists f. f \in X \rightarrow_M \mathbb{B} \wedge \mathcal{P}((x, y) \text{ in } \mathcal{D}. f \ x \neq y) = 0) \wedge$

$\mathcal{P}(s \text{ in } \Pi_M i \in \{..<m\}. \mathcal{D}. \mathcal{P}((x, y) \text{ in } \mathcal{D}. A \ s \ x \neq y) > 1 / 8) \geq 1 / 7$

# Lemma for the Proof of No-Free-Lunch Theorem

The following is used in the proof of the theorem.

## Lemma

$$A = \{x_0, y_0, x_1, y_1, \dots, x_n, y_n\}, \quad \implies \sum_{x \in A} f(x) = k * \frac{|A|}{2}$$
$$\forall i \leq n, f(x_i) + f(y_i) = k$$

**lemma** *sum-of-const-pairs*:

**fixes**  $k :: \text{real}$

**assumes** *finite A*

**and**  $\text{fst} \text{ ' } B \cup \text{snd} \text{ ' } B = A$  **and**  $\text{fst} \text{ ' } B \cap \text{snd} \text{ ' } B = \{\}$

**and** *inj-on fst B* **and** *inj-on snd B*

**and**  $\bigwedge x y. (x, y) \in B \implies f x + f y = k$

**shows**  $(\sum_{x \in A}. f x) = k * \text{real} (\text{card } A) / 2$

Note:

$$B = \{(x_0, y_0), \dots, (x_n, y_n)\}$$

Shown by induction on  $A$ .

# Conclusion

## No-Free-Lunch Theorem for ML in Isabelle/HOL

- Outline of the theorem
- Formalization in Isabelle/HOL
  - How to denote learning algorithms in Isabelle/HOL
  - A lemma used in the proof